

From: [Moody, Dustin \(Fed\)](#)
To: (b) (6)
Subject: RE: a very large Galois group, so that the number field is very far from having automorphisms.
Date: Thursday, April 26, 2018 2:47:00 PM

No.

On page 14 of their submission:

Current NTRU Classic specifications such as [32] prohibit m that have an unusually small number of 0's or 1's or -1's. For random m , this prohibition applies with probability $<2^{-10}$, and in case of failure the sender can try encoding the plaintext as a new m , but this is problematic for applications with hard real-time requirements. The reason for this prohibition is that NTRU Classic gives the attacker an "evaluate at 1" homomorphism from R/q to Z/q , leaking $m(1)$. The attacker scans many ciphertexts to find an occasional ciphertext where the value $m(1)$ is particularly far from 0; this value constrains the search space for the corresponding m by enough bits to raise security concerns. In NTRU Prime, R/q is a field, so this type of leak cannot occur.

The map $\phi: R/q \rightarrow Z/q$ is a ring homomorphism defined by evaluating at 1. The image of the map will be isomorphic to $(R/q) / (\ker \phi)$. We will be looking in the image of ϕ to see where $c(1)=m(1)$ is large. But if R/q is a field, then $\ker \phi = 1$ or $\ker \phi = R/q$. We know the kernel has more than one element, since any message with an equal number of 1s and -1s will be in the kernel. Hence, we must have the kernel of ϕ is all of R/q . Then this means the evaluation at 1 map is trivially the map which sends anything to 1, and reveals no information.

From: Quynh Dang (b) (6)
Sent: Thursday, April 26, 2018 2:40 PM
To: Moody, Dustin (Fed) <dustin.moody@nist.gov>
Subject: Re: a very large Galois group, so that the number field is very far from having automorphisms.

Are you coming in tomorrow ?

On Thu, Apr 26, 2018 at 2:37 PM, Moody, Dustin (Fed) <dustin.moody@nist.gov> wrote:

This seems to be true for NTRUprime. I don't see where they claim the attack doesn't work.

From: Quynh Dang (b) (6)
Sent: Thursday, April 26, 2018 2:30 PM
To: Moody, Dustin (Fed) <dustin.moody@nist.gov>
Subject: Re: a very large Galois group, so that the number field is very far from having automorphisms.

Here is the attack: Ciphertext : $c = rh + m$.

Number of 1s = number of -1s in r , so $r(1) = 0$ which implies $c(1) = r(1)h(1) + m(1) = m(1)$ which reveals information about m . If $c(1)$ is a huge positive number which means there are way more 1s than -1s which means that in m there are way more 1s than -1s: this gives information about m .

Quynh.

On Thu, Apr 26, 2018 at 2:25 PM, Moody, Dustin (Fed) <dustin.moody@nist.gov> wrote:

Write out the attack. Explain it to me....

From: Quynh Dang (b) (6)

Sent: Thursday, April 26, 2018 2:18 PM

To: Moody, Dustin (Fed) <dustin.moody@nist.gov>

Subject: Re: a very large Galois group, so that the number field is very far from having automorphisms.

Why not having subfield or subring stops that attack ?

Quynh.

On Thu, Apr 26, 2018 at 2:12 PM, Moody, Dustin (Fed) <dustin.moody@nist.gov> wrote:

That it doesn't have subrings (i.e. subfields), except the trivial ones.

From: Quynh Dang (b) (6)

Sent: Thursday, April 26, 2018 2:11 PM

To: Moody, Dustin (Fed) <dustin.moody@nist.gov>

Subject: Re: a very large Galois group, so that the number field is very far from having automorphisms.

So, what actually stops the attack in NTRU prime ?

On Thu, Apr 26, 2018 at 2:09 PM, Moody, Dustin (Fed) <dustin.moody@nist.gov> wrote:

Yes, it works on fields and rings. But it involves a subring....

From: Quynh Dang (b) (6)

Sent: Thursday, April 26, 2018 2:03 PM

To: Moody, Dustin (Fed) <dustin.moody@nist.gov>

Subject: Re: a very large Galois group, so that the number field is very far from having automorphisms.

Can you correct me below Dustin ?

On Thu, Apr 26, 2018 at 9:49 AM, Quynh Dang (b) (6) wrote:

Thank you Dustin.

Below is my understanding of the attack (wrong understanding).

Ciphertext : $c = rh + m$. Number of 1s = number of -1s in r , so $r(1) = 0$ which implies $c(1) = r(1)h(1) + m(1) = m(1)$

So, my wrong understanding is that the attack works for rings or fields.

Quynh.

On Thu, Apr 26, 2018 at 9:44 AM, Moody, Dustin (Fed) <dustin.moody@nist.gov> wrote:

Quynh,

I don't understand the statement about having a very large Galois group means the number field is very far from having automorphisms. By definition, the Galois group elements are automorphisms. So a large Galois group would mean a lot of automorphisms. I've read the blog, but I still can't make sense of it.

I think that a field blocks the evaluation at 1 attacks because the attack works with a subring. For a field, the subring is either the entire field or just $\{1\}$, which isn't helpful. By the way $\phi_n(x) = (x^n - 1)/(x - 1) = x^{n-1} + x^{n-2} + \dots + x + 1$. This will be irreducible if n is prime. The fact that it is irreducible means when we do $\mathbb{Q}[x] / \phi_n(x)$ we get a field.

Dustin

From: Quynh Dang (b) (6)

Sent: Thursday, April 26, 2018 8:59 AM

To: Moody, Dustin (Fed) <dustin.moody@nist.gov>

Subject: a very large Galois group, so that the number field is very far from having automorphisms.

Hi Dustin,

On a Dan's blog article: <https://blog.cr.yt.to/20140213-ideal.html>, he said that " and uses an irreducible polynomial $x^p - x - 1$ with a very large Galois group, so that the number field is very far from having automorphisms. " .

Why is this harder to find automorphisms if the Galois group is large ?

Why R/q (defined in NTRU prime) (a field instead of a ring) avoids evaluation at $m(1)$ attack? The attack seems to work as long as the number of -1 and 1 coefficients are known in r (I think my understanding for the attack is wrong here) because Tanga claims that replacing $X^N - 1$ in the original NTRU with $(X^N - 1)/(x - 1)$ to avoid the attack.

If the claim is correct, my impression is that $(X^N - 1)/(x - 1)$ is irreducible (I don't know this is true or not). If this is true, why does it being irreducible avoids the attack?

Thank you!
Quynh.